

High-Dimensional Adaptive Sparse Polynomial Interpolation and Applications to Parametric PDEs

Abdellah Chkifa · Albert Cohen ·
Christoph Schwab

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Abstract We consider the problem of Lagrange polynomial interpolation in high or countably infinite dimension, motivated by the fast computation of solutions to partial differential equations (PDEs) depending on a possibly large number of parameters which result from the application of generalised polynomial chaos discretisations to random and stochastic PDEs. In such applications there is a substantial advantage in considering polynomial spaces that are sparse and anisotropic with respect to the different parametric variables. In an adaptive context, the polynomial space is enriched at different stages of the computation. In this paper, we study an interpolation technique in which the sample set is incremented as the polynomial dimension increases, leading therefore to a minimal amount of PDE solving. This construction is based on the standard principle of tensorisation of a one-dimensional interpolation scheme and sparsification. We derive bounds on the Lebesgue constants for this interpolation pro-

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A. Chkifa · A. Cohen

Laboratoire Jacques-Louis Lions, UPMC Univ Paris 06, UMR 7598, 75005, Paris, France

A. Chkifa

e-mail: chkifa@ann.jussieu.fr

A. Cohen

e-mail: cohen@ann.jussieu.fr

A. Chkifa · A. Cohen

Laboratoire Jacques-Louis Lions, CNRS, UMR 7598, 75005, Paris, France

C. Schwab (✉)

Seminar for Applied Mathematics, ETH Zürich, 8092 Zürich, Switzerland

e-mail: schwab@math.ethz.ch

cess in terms of their univariate counterpart. For a class of model elliptic parametric PDE's, we have shown in Chkifa et al. (Modél. Math. Anal. Numér. 47(1):253–280, 2013) that certain polynomial approximations based on Taylor expansions converge in terms of the polynomial dimension with an algebraic rate that is robust with respect to the parametric dimension. We show that this rate is preserved when using our interpolation algorithm. We also propose a greedy algorithm for the adaptive selection of the polynomial spaces based on our interpolation scheme, and illustrate its performance both on scalar valued functions and on parametric elliptic PDE's.

Keywords Parametric and stochastic PDEs · Sparse polynomial approximation · High-dimensional problems · Adaptive algorithms · Lebesgue constants · Leja points

Mathematics Subject Classification 65N35 · 65L10 · 35J25 · 65D05

1 Introduction

In recent years, various strategies have been proposed for the numerical treatment of parametric partial differential equations [1, 3–5, 7, 11–13, 17, 18, 21–24, 27, 28, 30, 31]. Such equations have the general form

$$\mathcal{D}(u, \mathbf{y}) = 0, \quad (1.1)$$

where $u \mapsto \mathcal{D}(u, \mathbf{y})$ is a partial differential operator that depends on d parameters represented by the vector $\mathbf{y} = (y_1, \dots, y_d) \in \mathcal{P} \subset \mathbb{R}^d$. Assuming well-posedness of the problem in some Banach space V , the solution map,

$$\mathbf{y} \mapsto u(\mathbf{y}), \quad (1.2)$$

is defined from the parameter domain \mathcal{P} to the solution space V .

Parametric problems of this type arise in stochastic and deterministic modelling, depending on the nature of the parameters y_j , which may either be random or deterministic variables. In both settings, the main computational challenge is to approximate the entire solution map $\mathbf{y} \mapsto u(\mathbf{y})$ up to a prescribed accuracy, with reasonable computational cost. This task is particularly difficult when the number of involved parameters, d , is large due to the *curse of dimensionality*. This notion was coined by Bellmann in the early 1960s and refers to an exponential growth of the computational work for reaching a given target accuracy as the dimension d of the parameter space increases. In certain instances, the number of parameters may even be countably infinite, meaning that $\mathbf{y} = (y_j)_{j \geq 1}$.

High order polynomial approximation methods such as studied in [3, 5, 11–13, 18, 22, 23, 27, 30, 31], build approximations of the form

$$u_\Lambda(\mathbf{y}) = \sum_{\mathbf{v} \in \Lambda} u_{\mathbf{v}} \mathbf{y}^{\mathbf{v}}, \quad (1.3)$$

where $\Lambda \subset \mathcal{F}$ is a finite set of (multi-)indices $\mathbf{v} = (v_j)_{j \geq 1} \in \mathcal{F}$ and $\mathbf{y}^{\mathbf{v}} = \prod_{j \geq 1} y_j^{v_j}$. In the finite-dimensional setting $d < \infty$, the index set \mathcal{F} coincides with \mathbb{N}_0^d (here

and throughout, we denote by $\mathbb{N} = \{1, 2, 3, \dots\}$ the set of natural numbers and by $\mathbb{N}_0 = \mathbb{N} \cup \{0\}$. In the infinite-dimensional setting, \mathcal{F} denotes the (countable) set of all sequences of nonnegative integers which are *finitely supported* (i.e. those sequences for which only finitely many terms are non-zero). Note that the polynomial coefficients u_v are functions in V , and therefore the construction of u_Λ requires in principle the computation of $\#(\Lambda)$ such functions. This means that u_Λ is picked in the space

$$V_\Lambda = \text{Span} \left\{ \sum_{v \in \Lambda} v_v \mathbf{y}^v : v_v \in V \right\}. \quad (1.4)$$

We remark that

$$V_\Lambda := \mathbb{P}_\Lambda \otimes V, \quad (1.5)$$

where

$$\mathbb{P}_\Lambda := \text{Span} \{ \mathbf{y}^v : v \in \Lambda \} \quad (1.6)$$

denotes the polynomial space associated with the index set Λ and with coefficients in \mathbb{R} .

Polynomial approximation is well known to be effective when the solution map has some smoothness. In certain instances, it can even provably break the curse of dimensionality, in the sense that an algebraic convergence rate with respect to $\#(\Lambda)$ can be established even for functions of countably many parameters $d = \infty$. Such results were proved in [11–13] for the model parametric elliptic equation

$$-\text{div}(a \nabla u) = f \quad \text{in } D \subset \mathbb{R}^m, \quad u = 0 \text{ on } \partial D, \quad (1.7)$$

where $f \in H^{-1}(D)$ and $a(x, \mathbf{y}) := \bar{a}(x) + \sum_{j \geq 1} y_j \psi_j(x)$, with the functions ψ_j and \bar{a} in $L^\infty(D)$, and for the parameter domain $\mathcal{P} := [-1, 1]^\mathbb{N}$. More precisely, it was proved in [12, 13] that if a satisfies the *uniform ellipticity assumption*,

$$0 < r \leq a(x, \mathbf{y}) \leq R < \infty, \quad x \in D, \quad \mathbf{y} \in \mathcal{P}, \quad (1.8)$$

(UEA(r, R) for short) and if for some $0 < p < 1$, one has $(\|\psi_j\|_{L^\infty(D)})_{j \geq 1} \in \ell^p(\mathbb{N})$, then there exists a nested sequence

$$\Lambda_1 \subset \Lambda_2 \subset \dots \subset \Lambda_n \subset \dots \subset \mathcal{F}, \quad (1.9)$$

with $\#(\Lambda_n) = n$ which satisfies the uniform convergence bound

$$\inf_{v \in V_{\Lambda_n}} \sup_{\mathbf{y} \in \mathcal{P}} \|u(\mathbf{y}) - v(\mathbf{y})\|_V \leq C n^{-s}, \quad s := \frac{1}{p} - 1. \quad (1.10)$$

It was also shown in [12, 13] that there exists a sequence of sets Λ_n which ensures the mean square error convergence bound

$$\inf_{v \in V_{\Lambda_n}} \|u - v\|_{L^2(\mathcal{P}, V)} \leq C n^{-\bar{s}}, \quad \bar{s} := \frac{1}{p} - \frac{1}{2}, \quad (1.11)$$

where L^2 is defined with respect to the countable product of uniform measures on P . We point out that the results in this paper apply to a wide class of parametric problems. Problem (1.7) is chosen for the sake of illustration only; regularity estimates for the parametric solution analogous to those used in the present paper have been obtained for nonlinear, elliptic PDE in [20], for nonlinear, initial value ODEs in [21], and for parametric, parabolic evolution problems in [22], and for parametric wave equations in [23]. We also note that regularity for elliptic systems with operators depending on the parameter sequence \mathbf{y} have been considered in [29]. To minimise technicalities in presenting our results, we work throughout with (1.7), but investigate in the numerical experiments in Sect. 5 the performance of the adaptive interpolation algorithm also in other settings.

The results in [12, 13] were centred on the existence of such sequences of sets Λ_n rather than on their practical construction. The construction of sequences of (not necessarily optimal) sets Λ_n which achieve the convergence rates (1.10), (1.11), and therefore of the polynomial spaces \mathbb{P}_{Λ_n} , is critical in the design of algorithms for high-dimensional approximation. Sequences of quasioptimal sets Λ_n which give the rates (1.10), (1.11) with possibly a suboptimal constant $C > 0$ can either be derived from a priori estimates in [5, 12, 13] or by an adaptive search [11, 18]. The resulting spaces \mathbb{P}_{Λ_n} typically differ from the standard multivariate polynomial spaces \mathbb{P}_k of fixed total degree, and also from isotropic or anisotropic sparse grid polynomial spaces studied in [30, 31], although we sometimes refer to them as sparse polynomial spaces.

Given a finite index set Λ , several strategies can be used to compute $u_\Lambda \in V_\Lambda$:

- (i) Taylor expansions [11] can be recursively computed in the case of problems with affine parameter dependence such as (1.7). Adaptive methods based on such expansions have been proved to converge uniformly in $u \in \mathcal{P}$ with the same rate as in (1.10).
- (ii) Projection methods [1, 4, 5, 12, 18, 22, 23] produce near best approximations in V_Λ for the metric $L^2(\mathcal{P}, \rho)$ where ρ is a chosen measure in the parameter space. In addition, in the Galerkin framework, it is possible to use techniques of a-posteriori analysis in order to adaptively build the sequence of index sets $(\Lambda_n)_{n \geq 1}$. This approach was developed in [18] for the problem (1.7), and proved to converge with the same rate as in (1.11).
- (iii) Collocation methods [3, 5, 27, 30, 31] produce a polynomial approximation in V_Λ based on the data of particular solution instances $u(\mathbf{y}_i)$ for some chosen values $\mathbf{y}_1, \dots, \mathbf{y}_k$ of the parameter vector. One significant advantage of this approach is that it is *non-intrusive*: the sampling $u(\mathbf{y}_i)$ can be computed by any given numerical solver for the problem (1.1) and the polynomial approximation is built from these solutions by numerical techniques similar to those employed for scalar valued maps such as interpolation or least-square regression. On the other hand, the theoretical analysis of collocation methods is less satisfactory in the sense that convergence rates similar to (1.10) and (1.11) do not seem to have been established for such methods. This is in part due to the difficulty to control the stability of interpolation operators in arbitrary high dimension. In addition, adaptive methods for building the sets Λ_n have not been much developed in the collocation framework.

The objective of this paper is to propose and study a collocation method based on a high-dimensional interpolation process that can naturally be coupled with an adaptive selection of the polynomial spaces. We construct an interpolation operator I_Λ that maps real or complex valued functions defined on \mathcal{P} into \mathbb{P}_Λ . A standard vectorisation technique yields a similar operator that maps V -valued functions defined on \mathcal{P} into V_Λ . We assume that the parameter domain has a tensorised form,

$$\mathcal{P} = P^d \quad (1.12)$$

(with the understanding that \mathcal{P} denotes the countable Cartesian product when $d = \infty$), where P is a closed and bounded coordinate domain (typically, a bounded interval in \mathbb{R} or disk in \mathbb{C}). We impose a constraint on the index sets Λ that are considered, described by the following definition.

Definition 1.1 A nonempty set $\Lambda \subset \mathcal{F}$ is called monotone if

$$\nu \in \Lambda \quad \text{and} \quad \mu \leq \nu \Rightarrow \mu \in \Lambda, \quad (1.13)$$

where $\mu \leq \nu$ means that $\mu_j \leq \nu_j$ for all j .

Considering only polynomial spaces \mathbb{P}_Λ associated to such sets is very natural. In particular, monotonicity of the set Λ allows one to replace the monomials \mathbf{y}^ν in the definition of the spaces \mathbb{P}_Λ by any other tensorised basis $L_\nu(\mathbf{y}) = \prod_{j \geq 1} L_{\nu_j}(y_j)$ where $L_0 \equiv 1$ and $L_k(y_j)$ has degree exactly equal to k with respect to the coordinate y_j (for example, Legendre polynomials). Let us mention that polynomial spaces \mathbb{P}_Λ associated to monotone index sets Λ have, in fact, been introduced earlier in various contexts; we refer to [15, 25, 26] and the references therein and to Remark 2.2 ahead.

Importantly, it has been shown in [11] that the sets Λ_n achieving the convergence rate (1.10) for the problem (1.7) can be chosen from the restricted class of *monotone subsets of \mathcal{F}* . While we develop, as in [11], the algorithms and theory for (1.7), we hasten to add that *all results and algorithms presented in the present paper apply, without any modifications, to the adaptive numerical solution of more general parametric equations*: all that is required is bounded invertibility of the parametric equation for all instances of the parameter sequence and a characterisation of the parametric solution families' dependence on the parameters in the sequence. Such characterisations seem to hold for broad classes of parametric problems (we refer to [20–23] for details).

This paper is organised as follows. In Sect. 2, we build the operator I_Λ and its associated grid Γ_Λ , which we prove to be unisolvent for \mathbb{P}_Λ , for general monotone sets Λ . This construction is based on a univariate sequence of points $(z_k)_{k \geq 0}$ and a standard tensorisation and sparsification technique, originally due to Smolyak [32]. The main feature of this process is the inherent nested structure the grids, which is well adapted to an adaptive construction of the index set: the enrichment of Λ by one index is reflected by the enrichment of Γ_Λ by one point. The amount of computation is therefore minimised since all previously computed solution instances are used.

In Sect. 3, we study the stability of the interpolation operator I_Λ . In particular we establish bounds on the Lebesgue constant which only depends on the cardinality

of the set Λ (not on its shape or on the parametric dimension). These bounds grow algebraically with $\#(\Lambda)$, provided that bounds that grow algebraically with k are available for the Lebesgue constants associated to the sections $\{z_0, \dots, z_k\}$ of the sequence $(z_k)_{k \geq 0} \subset P$. Such univariate results have recently been obtained in [8–10] for particular choices of sequences.

Combining the approximation estimate (1.10) together with a bound of the form $(\#(\Lambda))^b$ for the Lebesgue constant, one expects that the interpolation $I_{\Lambda_n} u$ converges towards u with a deteriorated rate $n^{-(s-b)}$. In Sect. 4, we show by a different error analysis that, under the same assumption for the elliptic equation (1.7), one can construct a sequence $(\Lambda_n)_{n \geq 0}$ such that the interpolation converges towards u with the optimal rate n^{-s} . We present in Sect. 5 numerical results that illustrate the performance of our adaptive interpolation scheme.

2 Sparse Polynomial Interpolation

2.1 Properties of Monotone Sets

For $i \in \mathbb{N}$, we denote by $e_i = (\delta_{ij})_{j \geq 1}$ the i th Kronecker sequence. A monotone set Λ can equivalently be defined by the property

$$v \in \Lambda \quad \text{and} \quad v_i \neq 0 \Rightarrow v - e_i \in \Lambda. \quad (2.1)$$

Any monotone set contains the null multi-index $(0, 0, \dots)$, which we will denote by $0_{\mathcal{F}}$.

In addition to the partial order relation $\mu \leq v$ on \mathcal{F} , we say that $\mu < v$ if and only if $\mu \leq v$ but $\mu \neq v$ and $\mu \not\leq v$ if and only if $\mu_j > v_j$ for some $j \geq 1$. We say that an index v is *maximal* in a set $\Lambda \subset \mathcal{F}$ if and only if there is no $\mu \in \Lambda$ satisfying $v < \mu$. Any finite set has at least one maximal element. If Λ is monotone and if v is maximal in Λ , then $\Lambda \setminus \{v\}$ is monotone. Conversely, if $\tilde{\Lambda} \subset \Lambda$ are two monotone sets that differ by one element v , then v is maximal in Λ .

It follows that if $(\Lambda_n)_{n \geq 1}$ is a nested sequence of monotone sets with $\#(\Lambda_n) = n$, there exists a unique sequence of indices $(v^n)_{n \geq 1} \in \mathcal{F}^{\mathbb{N}}$ with $v^1 = 0_{\mathcal{F}}$ and such that

$$\Lambda_n = \{v^1, \dots, v^n\},$$

for all $n \geq 1$, with $v^1 = 0_{\mathcal{F}}$ and v^n maximal in Λ_n . Particular examples of monotone sets are the rectangles \mathcal{R}_v defined for any $v \in \mathcal{F}$ by

$$\mathcal{R}_v := \{\mu \in \mathcal{F} : \mu \leq v\}. \quad (2.2)$$

The only maximal element of a rectangle \mathcal{R}_v is v . In general, any finite monotone set $\Lambda \subset \mathcal{F}$ is completely determined by its maximal elements according to

$$\Lambda = \bigcup_{\substack{v \in \Lambda \\ v \text{ maximal}}} \mathcal{R}_v. \quad (2.3)$$

Conversely, given μ^1, \dots, μ^n , n multi-indices such that $i \neq j \Rightarrow \mu^i \not\leq \mu^j$, then $\bigcup_{j=1}^n \mathcal{R}_{\mu^j}$ is the only monotone set whose maximal elements are exactly the μ^j .

2.2 Univariate Interpolation and Tensorisation

Let $(z_k)_{k \geq 0}$ be a sequence of mutually distinct points in P (see (1.12)). We denote by I_k the univariate polynomial interpolation operator associated with the section $\{z_0, \dots, z_k\}$. This operator acts on a real or complex valued function g defined on P according to

$$I_k g := \sum_{i=0}^k g(z_i) l_i^k, \quad (2.4)$$

where

$$l_j^k(y) := \prod_{\substack{i=0 \\ i \neq j}}^k \frac{y - z_i}{z_j - z_i} \quad (2.5)$$

are the Lagrange polynomials associated with $\{z_0, \dots, z_k\}$. The possible dependence of the l_j^k on the interpolation points $\{z_0, \dots, z_k\}$ is expressed by the superscript k in (2.5). For $k \geq 0$, we introduce the difference operator

$$\Delta_k := I_k - I_{k-1}, \quad (2.6)$$

with the convention that I_{-1} is the null operator. Therefore, $\Delta_0 = I_0$ is the operator that maps g to the constant polynomial with value $g(z_0)$. With such notation, we write for any $n \in \mathbb{N}$

$$I_n = \sum_{k=0}^n \Delta_k. \quad (2.7)$$

Introducing the hierarchical polynomials of degree k associated to the sequence $(z_k)_{k \geq 0}$

$$h_k(y) := \prod_{j=0}^{k-1} \frac{y - z_j}{z_k - z_j}, \quad k > 0 \quad \text{and} \quad h_0(y) = 1, \quad (2.8)$$

one easily checks that

$$\Delta_k g = \alpha_k(g) h_k, \quad \alpha_k(g) := g(z_k) - I_{k-1} g(z_k), \quad (2.9)$$

and therefore we obtain the representation

$$I_n g = \sum_{k=0}^n \alpha_k(g) h_k, \quad (2.10)$$

which may be viewed as a re-expression of the standard Newton form. To any multi-index $v \in \mathcal{F}$, we associate the multivariate point

$$\mathbf{z}_v = (z_{v_j})_{j \geq 1} \in \mathcal{P}, \quad (2.11)$$

and the tensorised hierarchical function

$$\mathbf{H}_v(\mathbf{y}) = \prod_{j \geq 1} h_{v_j}(y_j), \quad (2.12)$$

and the tensorised multivariate operators

$$I_v := \bigotimes_{j \geq 1} I_{v_j} \quad \text{and} \quad \Delta_v := \bigotimes_{j \geq 1} \Delta_{v_j}. \quad (2.13)$$

The above tensorisation can be defined inductively: for a continuous and bounded, real or complex valued function g defined on \mathcal{P} ,

- If $v = 0_{\mathcal{F}}$, then $\Delta_v g = I_v g$ is the constant polynomial with value $g(\mathbf{z}_{0_{\mathcal{F}}})$.
- If $v \neq 0_{\mathcal{F}}$, then

$$I_v g = I_{v_1}(t \mapsto I_{\hat{v}} g_t) \quad \text{and} \quad \Delta_v g = \Delta_{v_1}(t \mapsto \Delta_{\hat{v}} g_t), \quad (2.14)$$

where $\hat{v} := (v_2, v_3, \dots)$ and where, for $t \in P$, the function g_t is defined on $P^{\mathbb{N}}$ by $g_t(\hat{\mathbf{y}}) = g(\mathbf{y})$, $\mathbf{y} := (t, \hat{\mathbf{y}})$.

In the case that $1 < d < \infty$, the function g_t is defined in a similar way but on P^{d-1} . In this case, the above induction terminates after exactly d steps. In the case $d = \infty$, the induction terminates after a finite number of steps since for any $v \in \mathcal{F}$, applying the operation $v \mapsto \hat{v}$ sufficiently many times to v leads to the null multi-index $0_{\mathcal{F}}$. We also observe that for every $v \in \mathcal{F}$

$$I_v = \bigotimes_{j \geq 1} \left(\sum_{k=0}^{v_j} \Delta_k \right) = \sum_{\mu \leq v} \Delta_{\mu}. \quad (2.15)$$

It is easily seen that I_v is the interpolation operator for the tensor product polynomial space $\mathbb{P}_{\mathcal{R}_v} = \bigotimes_{j \geq 1} \mathbb{P}_{v_j}$ associated to the tensorised grid $\Gamma_{\mathcal{R}_v} := \{\mathbf{z}_v : v \in \mathcal{R}_v\}$.

2.3 The Sparse Interpolation Operator

We are now ready to define the interpolation operator for polynomial spaces associated to monotone spaces and its corresponding unisolvent grid, following the technique of Smolyak [32]: for any monotone set $\Lambda \subset \mathcal{F}$, we define

$$I_{\Lambda} := \sum_{v \in \Lambda} \Delta_v, \quad (2.16)$$

and

$$\Gamma_{\Lambda} := \{\mathbf{z}_v : v \in \Lambda\}. \quad (2.17)$$

The monotone sets Λ might significantly differ from the sparse grid sets which are usually considered in the literature (see, e.g., [2, 6, 19, 32]). However, the arguments showing that I_{Λ} is the polynomial interpolation operator on \mathbb{P}_{Λ} associated to the grid Γ_{Λ} are very similar. For convenience of the reader, we give a precise statement of this result.

Theorem 2.1 *The grid Γ_Λ is unisolvent for \mathbb{P}_Λ and for any function g defined on $\overline{\mathcal{P}}$, the unique element in \mathbb{P}_Λ which agrees with g on Γ_Λ is given by $I_\Lambda g$.*

Proof It is readily seen that $I_\Lambda g$ belongs to \mathbb{P}_Λ for any function g defined on \mathcal{P} . Since we have

$$\#(\Gamma_\Lambda) = \#(\Lambda) = \dim(\mathbb{P}_\Lambda), \quad (2.18)$$

it suffices to verify that $I_\Lambda g$ agrees with g on Γ_Λ . For this we use (2.15) and the fact that Λ is a monotone set to write for any $v \in \Lambda$,

$$I_\Lambda = I_v + \sum_{\mu \in \Lambda, \mu \not\leq v} \Delta_\mu. \quad (2.19)$$

Since I_v is the interpolation operator for the space $\mathbb{P}_{\mathcal{R}_v}$ associated to the grid $\Gamma_{\mathcal{R}_v}$, we obviously have

$$\forall \mathbf{z}_v \in \Gamma_\Lambda: \quad I_v g(\mathbf{z}_v) = g(\mathbf{z}_v). \quad (2.20)$$

On the other hand if $\mu \not\leq v$, there exists a j such that $v_j < \mu_j$. Therefore, the univariate operator Δ_{μ_j} returns a polynomial which vanishes at z_{v_j} , and so $\Delta_\mu g$ vanishes at all points of j coordinate equal to z_{v_j} . In particular, therefore $\Delta_\mu g(\mathbf{z}_v) = 0$ and we have thus proved that $I_\Lambda g(\mathbf{z}_v) = g(\mathbf{z}_v)$. \square

Remark 2.2 The fact that Γ_Λ is unisolvent for the polynomial space \mathbb{P}_Λ when Λ is monotone appears to be known from early work on polynomial interpolation: see Chap. IV in the book [25], in which bivariate polynomials associated to monotone sets are referred to as “polynômes pleins”. This also appears as a particular case of the theory of the “least polynomial space” for interpolation of functions on general multivariate point sets, see in particular [15]. Here, polynomials associated to monotone sets Λ are referred to as “order closed polynomials” and are proved to be the least polynomial spaces for sets of the form Γ_Λ . Polynomial interpolation based on monotone sets has also recently been considered in [16].

Remark 2.3 One can generalise this construction in a straightforward way to tensorised domains of the more general form $\mathcal{P} = \prod_{j \geq 1} P_j$ with different univariate sequences $(z_k^j)_{k \geq 0}$ in each coordinate domain P_j . Another straightforward generalisation is when the univariate polynomial spaces \mathbb{P}_k are replaced by more general nested spaces S_k such that $\{z_0, \dots, z_k\}$ is unisolvent for S_k . Then Γ_Λ is unisolvent for the space

$$S_\Lambda = \bigoplus_{v \in \Lambda} \bigotimes_{j \geq 1} S_{v_j},$$

which generalises \mathbb{P}_Λ and the interpolation operator is defined in a similar manner as I_Λ . Sparse grid interpolation based on hierarchical finite element spaces are a particular instance of this generalisation.

Note that in our construction, any sequence $(z_k)_{k \geq 0}$ of mutually distinct points can be used. However, the choice of the univariate sequence is critical for the stability of the resulting multivariate interpolation operator I_A , expressed by the Lebesgue constant

$$\mathbb{L}_A := \sup_{g \in B(\mathcal{P}) - \{0\}} \frac{\|I_A g\|_{L^\infty(\mathcal{P})}}{\|g\|_{L^\infty(\mathcal{P})}}, \quad (2.21)$$

where

$$B(\mathcal{P}) := \left\{ u: \mathcal{P} \rightarrow \mathbb{C} : \sup_{y \in \mathcal{P}} |u(y)| < \infty \right\}, \quad (2.22)$$

denotes the space of bounded functions g on \mathcal{P} defined *everywhere* on \mathcal{P} .

Remark 2.4 When $d < \infty$ and $\mathcal{P} = P^d$, this is equivalent to taking the above supremum over $g \in C(\mathcal{P})$ the space of continuous functions on \mathcal{P} . When $d = \infty$ we use $B(\mathcal{P})$ in order to avoid a discussion of the topology that is used to define continuous functions of a countable number of variables. Similarly, we avoid measure theoretic considerations by using the convention that

$$\|g\|_{L^\infty(\mathcal{P})} := \sup_{y \in \mathcal{P}} |g(y)|, \quad (2.23)$$

for any $g \in B(\mathcal{P})$ (instead of $\text{esssup}_{y \in \mathcal{P}} |g(y)|$).

We are interested in choosing sequences $(z_k)_{k \geq 0}$ such that the univariate Lebesgue constants

$$\lambda_k = \max_{g \in C(P) - \{0\}} \frac{\|I_k g\|_{L^\infty(P)}}{\|g\|_{L^\infty(P)}}, \quad (2.24)$$

associated with the univariate operators I_k grow moderately with respect to k , since, as shown in Sect. 3, this allows us to derive estimates on \mathbb{L}_A .

A classical construction of such univariate sequences is by fixing $z_0 \in P$ and defining inductively

$$z_k := \text{Argmax}_{z \in P} \prod_{j=0}^{k-1} |z - z_j|. \quad (2.25)$$

Such $(z_k)_{k \geq 0}$ are called *Leja sequences* on P , and moderate, algebraic growth of the Lebesgue constants λ_k can be established in certain cases as we recall in Sect. 3. In addition, the choice of a Leja sequence for $(z_k)_{k \geq 0}$ has an interesting implication on the adaptive choice of the sets Λ_n as we explain in Sect. 2.5.

2.4 Hierarchical Computation of the Interpolation Operator

As explained in the introduction, we are interested in performing polynomial interpolation for a nested sequence of sets $(\Lambda_n)_{n \geq 1}$ with $n = \#(\Lambda_n)$. Accordingly the grids $(\Gamma_{\Lambda_n})_{n \geq 1} \subset P^{\mathbb{N}}$ are also nested. The sets $\Lambda_n \subset \mathcal{F}$ may either be fixed in advance, or adaptively chosen based on information gained at earlier computational steps.

In this setting, we have observed that each Λ_n can be viewed as the section $\{v^1, \dots, v^n\}$ of a sequence $(v^k)_{k \geq 1} \in \mathcal{F}^{\mathbb{N}}$. This observation leads to an efficient algorithm for the computation of $I_{\Lambda_n} g$ from $I_{\Lambda_{n-1}} g$ and of the value of g at the new point \mathbf{z}_{v^n} . Indeed, by tensorisation we observe that $\Delta_{v^n} g$ is a multiple of the tensorised hierarchical function \mathbf{H}_{v^n} defined in (2.12). Since $\mathbf{H}_{v^n}(\mathbf{z}_{v^n}) = 1$, it follows that

$$\begin{aligned} \Delta_{v^n} g &= \Delta_{v^n} g(\mathbf{z}_{v^n}) \mathbf{H}_{v^n} = (I_{\Lambda_n} g(\mathbf{z}_{v^n}) - I_{\Lambda_{n-1}} g(\mathbf{z}_{v^n})) \mathbf{H}_{v^n} \\ &= (g(\mathbf{z}_{v^n}) - I_{\Lambda_{n-1}} g(\mathbf{z}_{v^n})) \mathbf{H}_{v^n}, \end{aligned} \quad (2.26)$$

and therefore

$$I_{\Lambda_n} g = I_{\Lambda_{n-1}} g + (g(\mathbf{z}_{v^n}) - I_{\Lambda_{n-1}} g(\mathbf{z}_{v^n})) \mathbf{H}_{v^n}. \quad (2.27)$$

Consequently, the polynomials $I_{\Lambda_n} g$ are given by

$$I_{\Lambda_n} g = \sum_{k=0}^n g_{v^k} \mathbf{H}_{v^k}, \quad (2.28)$$

where the coefficients g_{v^k} are defined recursively by

$$\begin{aligned} g_{v^1} &= g(\mathbf{z}_0), \\ g_{v^{k+1}} &:= g(\mathbf{z}_{v^{k+1}}) - I_{\Lambda_k} g(\mathbf{z}_{v^{k+1}}) = g(\mathbf{z}_{v^{k+1}}) - \sum_{i=1}^k g_{v^i} \mathbf{H}_{v^i}(\mathbf{z}_{v^{k+1}}). \end{aligned} \quad (2.29)$$

Remark 2.5 In the sum that appears on the right side of (2.29), only the terms such that $v^i \leq v^{k+1}$ are non-zero. When evaluating the computational cost in the above operation, one should make the distinction between the cost of the evaluation of $g(\mathbf{z}_{v^{k+1}})$ and of computing the linear combination $\sum_{i=1}^k g_{v^i} \mathbf{H}_{v^i}(\mathbf{z}_{v^{k+1}})$. In instances where the evaluation of g requires running a heavy numerical code (for example when $g(\mathbf{y})$ is an output of the solution $u(\mathbf{y})$ to a parametric PDE), the first cost dominates the second one. It is also important to notice that only n evaluations of g have been used until the current step n .

Remark 2.6 The above algorithm is also efficient to construct the interpolant $I_{\Lambda} g$ for any given monotone set Λ . Indeed, by iteratively removing maximal elements, we see that any such set can be written as $\Lambda = \Lambda_k$ with $k := \#(\Lambda)$ and $(\Lambda_n)_{n \geq 1}$ a sequence of the above type. It is also easily checked that the coefficients g_{v^k} only depend on g and on the index v^k and are independent on the index set Λ : these coefficients can be viewed as the unique coordinates of g in the hierarchical basis $(\mathbf{H}_v)_{v \in \mathcal{F}}$. One should, however, be cautious when writing the expansion

$$g = \sum_{v \in \mathcal{F}} g_v \mathbf{H}_v, \quad (2.30)$$

since it may fail to converge for certain functions g regardless of the ordering of the summation. However, it will be proved to converge for functions that can be approximated sufficiently well by polynomials, based on the stability analysis of the interpolation operator which is the object of the next section.

2.5 Adaptive Selection of Polynomial Spaces

We now discuss the adaptive selection of a nested sequence $(\Lambda_n)_{n \geq 1}$. Let us begin with the following analogy: if $(\mathbf{H}_v)_{v \in \mathcal{F}}$ was an orthonormal basis of $L^2(\mathcal{P})$ then the choice of an index set Λ_n that minimises the L^2 error when truncating the expansion (2.30) would be the indices corresponding to the n largest $|g_v|$.

In our current setting however, $(\mathbf{H}_v)_{v \in \mathcal{F}}$ is not an orthonormal basis and we are rather interested in controlling the error in L^∞ . A first greedy strategy is to define Λ_n as the set of indices corresponding to the n largest contributions of (2.30) measured in the L^∞ metric, i.e. the n largest $a_v |g_v|$, where

$$a_v := \|\mathbf{H}_v\|_{L^\infty(\mathcal{P})} := \prod_{j \geq 1} \|h_{v_j}\|_{L^\infty(P)}. \quad (2.31)$$

This strategy obviously gives rise to a nested sequence $(\Lambda_n)_{n \geq 1}$, however, the sets Λ_n are not ensured to be monotone. In addition, it is not computationally feasible since finding the n largest contributions in (2.30) hints that we should have computed all contributions.

In order to correct these defects, we define for any monotone set Λ a set of neighbours

$$\mathcal{N}(\Lambda) := \{v \notin \Lambda : \mathcal{R}_v \in \Lambda \cup \{v\}\}, \quad (2.32)$$

or equivalently those $v \notin \Lambda$ such that $v - e_j \in \Lambda$ for all j such that $v_j \neq 0$. Then, a natural variant of the first strategy that leads to a nested sequence of monotone sets is the following greedy adaptive algorithm.

Adaptive Interpolation (AI) Algorithm:

- Start with $\Lambda_1 := \{0_{\mathcal{F}}\}$.
- Assuming that Λ_{n-1} has been computed, find

$$v^n := \operatorname{argmax}\{a_v |g_v| : v \in \mathcal{N}(\Lambda_{n-1})\}, \quad (2.33)$$

and define $\Lambda_n = \Lambda_{n-1} \cup \{v^n\}$.

Let us observe that since $\mathbf{H}_v(\mathbf{z}_v) = 1$, we obviously have $a_v \geq 1$. On the other hand, when $(z_k)_{k \geq 0}$ is a Leja sequence built according to (2.25), we obviously have $\max_{z \in P} |h_k(z)| = |h_k(z_k)| = 1$ and therefore

$$a_v = \mathbf{H}_v(\mathbf{z}_v) = 1. \quad (2.34)$$

In such a case, in view of (2.29), the greedy strategy (2.33) amounts to choosing the new index in $\mathcal{N}(\Lambda_{n-1})$ that maximises the interpolation error at the corresponding new grid point:

$$v^n := \operatorname{argmax}\{|g(\mathbf{z}_v) - I_{\Lambda_{n-1}}g(\mathbf{z}_v)| : v \in \mathcal{N}(\Lambda_{n-1})\}. \quad (2.35)$$

This greedy strategy has several defects. The first one is that it may simply fail to converge, even if there exist sequences $(\Lambda_n)_{n \geq 0}$ such that $I_{\Lambda_n}g$ converges to g

at a high rate. This is due to data oscillation that could return an artificially small interpolation error at the new grid point. Consider for example a two-dimensional function of the form

$$g(\mathbf{y}) = g(y_1, y_2) = g_1(y_1)g_2(y_2), \quad (2.36)$$

where g_1 and g_2 are non-polynomial smooth functions such that g_2 takes the same values at the points z_0 and z_1 . Then, the algorithm will select sets Λ_n that consist of the indices $v = (k, 0)$ for $k = 0, \dots, n-1$, since the interpolation error at the point $\mathbf{z}_{(k,1)} = (z_k, z_1)$ will vanish. Although this type of situation might be viewed as pathological, it reflects the fact that the algorithm might fail in its first steps to identify the significant variables. One way to avoid this is to impose that when all interpolation errors $|g(\mathbf{z}_v) - I_{\Lambda_{n-1}}g(\mathbf{z}_v)|$ for $v \in \mathcal{N}(\Lambda_n)$ are smaller than some prescribed tolerance $\varepsilon_n > 0$ (that is either fixed or tends to 0 as n grows), then the new index v^n is chosen arbitrarily from $\mathcal{N}(\Lambda_n)$.

The second defect is that in the infinite-dimensional framework $d = \infty$ the set of neighbours $\mathcal{N}(\Lambda)$ has infinite cardinality. One way to treat this defect is to modify its definition by setting

$$\mathcal{N}(\Lambda) := \{v \notin \Lambda : v_j = 0 \text{ if } j > j(\Lambda) + 1 \text{ and } \mathcal{R}_v \in \Lambda \cup \{v\}\}, \quad (2.37)$$

where $j(\Lambda) := \max\{j : v_j > 0 \text{ for some } v \in \Lambda\}$. This means that we can activate at most one new variable at each iteration step.

Even with such modifications, it is not clear to understand under which additional assumptions on g this adaptive, greedy selection procedure will pick sets $(\Lambda_n)_{n \geq 1}$ such that the interpolation $I_{\Lambda_n}g$ has a guaranteed convergence rate comparable to that of an optimal choice of sets such as, e.g. obtained from the best n -term approximation. We give in Sect. 5 several numerical examples that illustrate the good practical behaviour of this algorithm.

3 The Lebesgue Constant

The accuracy of the interpolation operator I_Λ can be related to the error of best polynomial approximation via the Lebesgue constant \mathbb{L}_Λ defined in (2.21), according to the classical inequality

$$\|g - I_\Lambda g\|_{L^\infty(\mathcal{P})} \leq (1 + \mathbb{L}_\Lambda) \inf_{Q \in \mathbb{P}_\Lambda} \|g - Q\|_{L^\infty(\mathcal{P})}. \quad (3.1)$$

A crude, yet useful, way to estimate \mathbb{L}_Λ is by using the triangle inequality which gives

$$L_\Lambda \leq \sum_{v \in \Lambda} \delta_v, \quad (3.2)$$

where we define for $v \in \mathcal{F}$

$$\delta_v := \sup_{g \in B(\mathcal{P}) - \{0\}} \frac{\|\Delta_v g\|_{L^\infty(\mathcal{P})}}{\|g\|_{L^\infty(\mathcal{P})}}, \quad (3.3)$$

where $B(\mathcal{P})$ is defined as in Sect. 2.3. It is readily seen that

$$\delta_v := \prod_{j \geq 1} \delta_{v_j}, \quad (3.4)$$

where

$$\delta_k := \sup_{g \in C(P) - \{0\}} \frac{\|\Delta_k g\|_{L^\infty(P)}}{\|g\|_{L^\infty(P)}} \leq \lambda_{k-1} + \lambda_k, \quad (3.5)$$

with λ_k the Lebesgue constant associated with the univariate operators I_k and with the convention that $\lambda_{-1} := 0$. Therefore

$$\mathbb{L}_\Lambda \leq \sum_{v \in \Lambda} \prod_{j \geq 1} (\lambda_{v_j} + \lambda_{v_{j-1}}). \quad (3.6)$$

Note that $\delta_0 = \lambda_0 = 1$, regardless of the choice of the sequence $(z_k)_{k \geq 0}$.

The bound (3.6) is of course crude, since we did not take advantage of the telescoping nature in the summation of the Δ_v . For instance, when $\Lambda = \mathcal{R}_v$, we have seen that $I_\Lambda = I_v = \bigotimes_{j \geq 1} I_{v_j}$, and in that case the exact value of the Lebesgue constant is given by the smaller value

$$\mathbb{L}_{\mathcal{R}_v} = \prod_{j \geq 1} \lambda_{v_j}. \quad (3.7)$$

Nevertheless, for general monotone sets Λ , we can use the bound (3.6) to study the behaviour of the Lebesgue constant \mathbb{L}_Λ as the number of interpolation points $\#(\Gamma_\Lambda) = \#(\Lambda)$ grows. The following result shows that when certain algebraic bounds are available for the λ_k , then similar algebraic bounds can be derived for \mathbb{L}_Λ in terms of $\#(\Lambda)$ regardless of the dimension d and of the shape of Λ .

Lemma 3.1 *If the Lebesgue constants λ_k satisfy*

$$\lambda_k \leq (k+1)^\theta, \quad k \geq 0 \quad (3.8)$$

for some $\theta \geq 1$, then the Lebesgue constants \mathbb{L}_Λ satisfy

$$\mathbb{L}_\Lambda \leq (\#(\Lambda))^{\theta+1} \quad (3.9)$$

for any monotone set Λ .

Proof For any $k \geq 0$, one has $\lambda_k + \lambda_{k-1} \leq (k+1)^\theta + k^\theta \leq (2k+1)(k+1)^{\theta-1}$, therefore, for $v \in \Lambda$

$$\begin{aligned} \prod_{j \geq 1} (\lambda_{v_j} + \lambda_{v_{j-1}}) &\leq \left(\prod_{j \geq 1} (v_j + 1) \right)^{\theta-1} \prod_{j \geq 1} (2v_j + 1) \\ &= (\#(\mathcal{R}_v))^{\theta-1} \prod_{j \geq 1} (2v_j + 1) \\ &\leq (\#(\Lambda))^{\theta-1} \prod_{j \geq 1} (2v_j + 1), \end{aligned}$$

where we have used $\mathcal{R}_v \subset \Lambda$ since Λ is monotone. To complete the proof, it remains to show that $\sigma(\Lambda) \leq (\#\Lambda)^2$, where

$$\sigma(\Lambda) := \sum_{v \in \Lambda} \prod_{j \geq 1} (2v_j + 1). \quad (3.10)$$

This is done using induction on $n_\Lambda := \#\Lambda$. When $n_\Lambda = 1$, then $\Lambda = \{0_{\mathcal{F}}\}$ and it is obviously true. Let $n \geq 1$ and Λ denote a monotone set with $n_\Lambda = n + 1$. Without loss of generality, we suppose that $\mu_1 \neq 0$ for some $\mu \in \Lambda$, and denote by $K \geq 1$ the maximal value attained by the coordinate v_1 when $v \in \Lambda$. For $0 \leq k \leq K$, we introduce the sections

$$\Lambda_k := \{\hat{v} = (v_2, v_3, \dots) : (k, \hat{v}) \in \Lambda\}. \quad (3.11)$$

By monotonicity of Λ , every Λ_k is monotone. Also, since $K \geq 1$ then $\#(\Lambda_k) < \#(\Lambda)$ for any k , so that the induction hypothesis implies

$$\sigma(\Lambda) = \sum_{k=0}^K (2k+1) \sigma(\Lambda_k) \leq \sum_{k=0}^K (2k+1) (\#(\Lambda_k))^2. \quad (3.12)$$

Also, we have

$$\Lambda_K \subset \dots \subset \Lambda_1 \subset \Lambda_0, \quad (3.13)$$

since for $k \geq 1$, $\mu \in \Lambda_k \Rightarrow (k, \mu) \in \Lambda \Rightarrow (k-1, \mu) \in \Lambda \Rightarrow \mu \in \Lambda_{k-1}$. We deduce

$$k(\#(\Lambda_k))^2 \leq \#(\Lambda_k) \#(\Lambda_0) + \dots + \#(\Lambda_k) \#(\Lambda_{k-1}), \quad (3.14)$$

and consequently

$$\begin{aligned} \sigma(\Lambda) &\leq \sum_{k=0}^K (\#(\Lambda_k))^2 + 2 \sum_{k=1}^K (\#(\Lambda_k) \#(\Lambda_0) + \dots + \#(\Lambda_k) \#(\Lambda_{k-1})) \\ &= \left(\sum_{k=0}^K \#(\Lambda_k) \right)^2 = (\#(\Lambda))^2, \end{aligned} \quad (3.15)$$

which concludes the proof. \square

Remark 3.2 In the case where $(z_k)_{k \geq 0}$ is a Leja sequence defined by (2.25) for some initial point $z_0 \in P$, the hierarchical polynomials h_k defined by (2.8) satisfy

$$|h_k(z)| \leq |h_k(z_k)| = 1, \quad z \in P. \quad (3.16)$$

Since, according to (2.9), we have

$$\Delta_k g = (g(z_k) - I_{k-1} g(z_k)) h_k, \quad (3.17)$$

it follows that

$$\delta_k \leq 1 + \lambda_{k-1}, \quad (3.18)$$

and

$$\mathbb{L}_\Lambda \leq \sum_{v \in \Lambda} \prod_{j \geq 1} (1 + \lambda_{v_j - 1}), \quad (3.19)$$

which are improvements over (3.5) and (3.19).

Let us observe that since $\lambda_{-1} = 0$ and $\lambda_0 = 1$, bounds of the form $\lambda_k \leq (k+1)^\theta$ can be established for some $\theta > 0$ provided that λ_k are bounded as $\mathcal{O}(k^b)$ for some $b > 0$. Such bounds have recently been obtained in [8–10], in the cases where P is either the complex unit disk $\{|z| \leq 1\}$ or the unit interval $[-1, 1]$.

For the complex unit disk, if $(z_k)_{k \geq 0}$ is the Leja sequence defined by (2.25) with $z_0 = 1$, it is proved in [8] that

$$\lambda_k = \mathcal{O}(k \log k)$$

and conjectured that

$$\lambda_k \leq k + 1, \quad k \geq 0. \quad (3.20)$$

The more precise bound

$$\lambda_k \leq 2(k + 1), \quad k \geq 1, \quad (3.21)$$

is established in [10]. Using the improvement (3.21), and using the same computation as in the proof of the above lemma, we obtain in that case that $\mathbb{L}_\Lambda \leq (\#(\Lambda))^2$.

For the unit interval $[-1, 1]$, and the sequence given by the projection of the above complex Leja sequence, an asymptotic bound into $\mathcal{O}(k^3 \log k)$ has been obtained in [9], and it has been improved to $3(k+1)^2 \log(k+1)$ in [10]. Note that projections of Leja sequences on the unit disk to $[-1, 1]$ are not Leja sequences on $[-1, 1]$.

4 Application to Parametric PDEs

4.1 Interpolation of Banach Valued Functions

We are interested in applying our interpolation process to the map $\mathbf{y} \mapsto u(\mathbf{y})$ defined by exact or approximate solution of the parametric PDE (1.1) for the given parameter \mathbf{y} . Therefore, we want to interpolate a function which is not real or complex valued, but instead takes its value in the solution space V . The generalisation of the interpolation operator I_Λ to this setting is straightforward: $I_\Lambda u$ is the unique function in V_Λ that coincides with u at the points \mathbf{z}_v for $v \in \Lambda$. As in the scalar case, it can be expanded according to

$$I_\Lambda u = \sum_{v \in \Lambda} u_v \mathbf{H}_v, \quad (4.1)$$

where the coefficients $u_v \in V$ can be computed in a recursive way similar to (2.29):

$$u_{v^1} = u(\mathbf{z}_0), \quad u_{v^{k+1}} = u(\mathbf{z}_{v^{k+1}}) - \sum_{i=1}^k u_{v^i} \mathbf{H}_{v^i}(\mathbf{z}_{v^{k+1}}), \quad (4.2)$$

where $\Lambda_n = \{v^1, \dots, v^n\}$, $n = 1, 2, \dots$, is a nested sequence of monotone sets. We are interested in the accuracy of the interpolant in the sense of the supremum error

$$\|u - I_\Lambda u\|_{L^\infty(\mathcal{P}, V)}, \quad (4.3)$$

where for a function v in the space $B(\mathcal{P}, V)$ of V valued bounded functions,

$$\|v\|_{L^\infty(\mathcal{P}, V)} := \sup_{\mathbf{y} \in \mathcal{P}} \|v(\mathbf{y})\|_V. \quad (4.4)$$

Remark 4.1 Similar to (2.23), we avoid measure theoretic considerations in the infinite-dimensional framework by using the exact supremum instead of the essential supremum in the definition of the $L^\infty(\mathcal{P}, V)$ norm. In practical instances of solution maps $\mathbf{y} \mapsto u(\mathbf{y})$ associated with parametric PDE's such as (1.7), it can be checked that the map is continuous over \mathcal{P} with respect to the Tychonoff topology (that is, the topology of pointwise convergence for bounded sequences) for which \mathcal{P} is compact, and therefore this supremum is attained.

The same reasoning as for interpolation of scalar valued functions shows that

$$\|u - I_\Lambda u\|_{L^\infty(\mathcal{P}, V)} \leq (1 + \mathbb{L}_\Lambda) \inf_{v \in V_\Lambda} \|u - v\|_{L^\infty(\mathcal{P}, V)}, \quad (4.5)$$

where \mathbb{L}_Λ is the Lebesgue constant associated to the interpolation operator I_Λ which was defined and studied in the previous section.

4.2 Convergence Rates for a Parametric, Elliptic Model Problem

As already explained in the introduction, for the model elliptic problem (1.7), one can establish convergence rates in $L^\infty(\mathcal{P}, V)$ and $L^2(\mathcal{P}, V)$, where $V = H_0^1(D)$, for polynomial approximations that are robust with respect to the parametric dimension. Here we focus on $L^\infty(\mathcal{P}, V)$. The results in [11–13] show that if the diffusion satisfies the uniform ellipticity assumption (1.8) and if for some $0 < p < 1$, one has $(\|\psi_j\|_{L^\infty(D)})_{j \geq 1} \in \ell^p(\mathbb{N})$, there exists a nested sequence $(\Lambda_n)_{n \geq 1}$ of monotone sets such that $\#(\Lambda_n) = n$ and

$$\inf_{v \in V_{\Lambda_n}} \|u - v\|_{L^\infty(\mathcal{P}, \Lambda)} \leq C n^{-s}, \quad s := \frac{1}{p} - 1. \quad (4.6)$$

In fact this rate of convergence was proved for specific approximations defined either by Taylor expansion

$$T_{\Lambda_n} u(\mathbf{y}) := \sum_{v \in \Lambda_n} t_v \mathbf{y}^v, \quad (4.7)$$

where $t_v := \frac{1}{v!} \partial_y^v u(\mathbf{y})|_{\mathbf{y}=0}$ with $v! := \prod_{j \geq 1} v_j!$ or by orthogonal projections

$$P_{\Lambda_n} u(\mathbf{y}) := \sum_{v \in \Lambda_n} c_v L_v(\mathbf{y}), \quad (4.8)$$

where the L_v are the tensorised Legendre polynomials in $[-1, 1]$, which are normalised with respect to the probability measure $dt/2$ and where c_v are the corresponding Legendre coefficients. The fact that we work in the infinite-dimensional case $d = \infty$ in (4.6) reveals that the convergence rate n^{-s} in (4.6) is robust with respect to the number of active variables. In this section, we work under the same assumptions on the diffusion coefficient and its expansion.

We now study the rate of convergence of $I_{\Lambda_n}u$ towards u . Combining (4.6) and (4.5), we obtain

$$\|u - I_{\Lambda_n}u\|_{L^\infty(\mathcal{P}, V)} \leq C(1 + \mathbb{L}_{\Lambda_n})n^{-s}. \quad (4.9)$$

We have seen in Sect. 4 that the Lebesgue constant can be controlled by a bound of the form

$$\mathbb{L}_{\Lambda_n} \leq n^b, \quad \text{where } b = \theta + 1, \quad (4.10)$$

when the univariate sequence $(z_k)_{k \geq 0}$ is chosen so that $\lambda_k \leq (k+1)^\theta$ for some $\theta > 0$ and we gave examples of such sequences. We thus obtain a convergence estimate of the form

$$\|u - I_{\Lambda_n}u\|_{L^\infty(\mathcal{P}, V)} \leq Cn^{-(s-b)}. \quad (4.11)$$

With this simple stability analysis (via the bound for the Lebesgue constant) plus consistency analysis (via the n -term approximation result), the convergence rate obtained in (4.11) is reduced by $b = \theta + 1$ compared to the (benchmark) n -term approximation rate s in (4.6).

We now present a more refined argument that shows that there exists a sequence $(\Lambda_n)_{n \geq 1} \subset \mathcal{F}$ of monotone sets such that $\#(\Lambda_n) = n$ and a constant $C > 0$ such that for all $n \in \mathbb{N}$ holds

$$\|u - I_{\Lambda_n}u\|_{L^\infty(\mathcal{P}, V)} \leq Cn^{-s}. \quad (4.12)$$

This analysis relates more directly the interpolation error with the Taylor coefficients of u according to the following result.

Lemma 4.2 Assume that $u(\mathbf{y}) = \sum_{v \in \mathcal{F}} t_v \mathbf{y}^v$ in the sense of unconditional convergence in $L^\infty(\mathcal{P}, V)$. If the univariate sequence $(z_k)_{k \geq 0}$ is chosen so that $\lambda_k \leq (k+1)^\theta$, then

$$\|(I - I_\Lambda)u\|_{L^\infty(\mathcal{P}, V)} \leq 2 \sum_{v \notin \Lambda} p_v(b) \|t_v\|_V, \quad (4.13)$$

for any finite monotone set Λ , where $b := \theta + 1$ and

$$p_v(b) := \prod_{j \geq 1} (1 + v_j)^b. \quad (4.14)$$

Proof The unconditional convergence of the Taylor series allows us to write

$$I_\Lambda u = I_\Lambda \left(\sum_{v \in \mathcal{F}} t_v \mathbf{y}^v \right) = \sum_{v \in \mathcal{F}} t_v I_\Lambda \mathbf{y}^v = \sum_{v \in \Lambda} t_v \mathbf{y}^v + \sum_{v \notin \Lambda} t_v I_{\Lambda \cap \mathcal{R}_v} \mathbf{y}^v. \quad (4.15)$$

Here, we used the fact that for monotone Λ we have $I_\Lambda \mathbf{y}^\nu = \mathbf{y}^\nu$ for every $\nu \in \Lambda$ and that $\Delta_\mu \mathbf{y}^\nu = 0$ for every $\mu \not\leq \nu$. Therefore

$$(I - I_\Lambda)u = \sum_{\nu \notin \Lambda} t_\nu (I - I_{\Lambda \cap \mathcal{R}_\nu}) \mathbf{y}^\nu, \quad (4.16)$$

where I stands for the identity operator. This implies

$$\|(I - I_\Lambda)u\|_{L^\infty(\mathcal{P}, V)} \leq \sum_{\nu \notin \Lambda} \|t_\nu\|_V (1 + \mathbb{L}_{\Lambda \cap \mathcal{R}_\nu}) \leq 2 \sum_{\nu \notin \Lambda} \|t_\nu\|_V \mathbb{L}_{\Lambda \cap \mathcal{R}_\nu}. \quad (4.17)$$

Since $\lambda_k \leq (k + 1)^\theta$, we obtain from Lemma 3.1 that

$$\mathbb{L}_{\Lambda \cap \mathcal{R}_\nu} \leq (\#(\Lambda \cap \mathcal{R}_\nu))^{\theta+1} \leq (\#(\mathcal{R}_\nu))^{\theta+1} = p_\nu(b), \quad (4.18)$$

which completes the proof. \square

Let us observe that the above lemma can be generalised to polynomial expansions other than Taylor series, for example expansions into Legendre polynomials, provided that unconditional convergence holds. We focus here on the Taylor series, which allows us to use the results in [13] that prove unconditional convergence under our assumptions on the diffusion coefficients, based on explicit bounds for the V -norms of Taylor coefficients. These bounds are obtained by an application of the Cauchy formula, on the holomorphic extension of the map $\mathbf{y} \mapsto u(\mathbf{y})$, and are of the form

$$\|t_\nu\|_V \leq C_\delta \prod_{j \geq 1} \rho_j^{-\nu_j} = C_\delta \rho^{-\nu}, \quad (4.19)$$

where $C_\delta := \frac{\|f\|_{V^*}}{\delta}$ for arbitrary $\delta > 0$ and $\rho := (\rho_j)_{j \geq 1}$ is any sequence of strictly positive numbers that satisfies (see (1.7))

$$\sum_{j \geq 1} \rho_j |\psi_j(x)| \leq \bar{a}(x) - \delta, \quad x \in D. \quad (4.20)$$

We call sequences ρ which satisfy (4.20) δ -admissible. We observe that δ has to be chosen smaller than the ellipticity constant r in order to ensure the existence of such sequences. In [13], (4.19) is used to prove the unconditional convergence of the Taylor series of u in $L^\infty(\mathcal{P}, V)$ and the convergence rate (4.6). In the next result, we use the same bounds to prove a summability result on the sequence $(p_\nu(b) \|t_\nu\|_V)_{\nu \in \mathcal{F}}$.

Theorem 4.3 *Assume that if the diffusion coefficient satisfies the uniform ellipticity assumption (1.8) and that for some $0 < p < 1$, one has $(\|\psi_j\|_{L^\infty(D)})_{j \geq 1} \in \ell^p(\mathbb{N})$. Then for any $b \geq 1$ and for $p_\nu(b)$ as in (4.14), the sequence $(p_\nu(b) \|t_\nu\|_V)_{\nu \in \mathcal{F}}$ belongs to $\ell^p(\mathcal{F})$.*

Proof We fix $B = e^b$. Since $p < 1$, $(\|\psi_j\|_{L^\infty(D)})_{j \geq 1} \in \ell^1(\mathbb{N})$ and we may choose $\kappa > 1$ and an integer $J_0 \geq 1$ such that

$$(\kappa - 1) \sum_{j \geq 1} \|\psi_j\|_{L^\infty(D)} \leq \frac{r}{6} \quad \text{and} \quad \sum_{j > J_0} \|\psi_j\|_{L^\infty(D)} \leq \frac{1}{B} \frac{r}{12}, \quad (4.21)$$

and we set $E := \{j: 1 \leq j \leq J_0\}$ and $F := \{j: j > J_0\}$. For each $v \in \mathcal{F}$, we denote by v_E and v_F the restrictions of v on E and F , and we define the sequence $\rho = \rho(v)$ by $\rho_j = \kappa$ if $j \leq J_0$, $\rho_j = 1$ if $j > J_0$ with $v_j = 0$, and

$$\rho_j = B + \frac{rv_j}{4|v_F|\|\psi_j\|_{L^\infty(D)}}, \quad j > J_0, \quad v_j \neq 0 \quad (4.22)$$

where $|v_F| = \sum_{j > J_0} v_j$. We claim that ρ is $\frac{r}{2}$ -admissible. Indeed, for almost every $x \in D$

$$\begin{aligned} \sum_{j \geq 1} \rho_j |\psi_j(x)| &\leq \kappa \sum_{j \leq J_0} |\psi_j(x)| + \frac{r}{4} \sum_{j > J_0} \frac{v_j}{|v_F|\|\psi_j\|_{L^\infty(D)}} |\psi_j(x)| + B \sum_{j > J_0} |\psi_j(x)| \\ &\leq \sum_{j \leq J_0} |\psi_j(x)| + \frac{r}{6} + \frac{r}{4} + \frac{r}{12} \\ &\leq \bar{a}(x) - r + \frac{r}{2} \\ &= \bar{a}(x) - \frac{r}{2} \end{aligned}$$

where we used the uniform ellipticity assumption (1.8). Therefore, by (4.19)

$$p_v(b)\|t_v\|_V \leq C_{\frac{r}{2}} \prod_{j \leq J_0} \frac{(1+v_j)^b}{\kappa^{v_j}} \prod_{j > J_0} \frac{(1+v_j)^b}{\rho_j^{v_j}}. \quad (4.23)$$

Since $\kappa > 1$, there exists $C_0 = C_0(b, \kappa) > 0$ such that $(1+n)^b \leq C_0(\frac{1+\kappa}{2})^n$ for any $n \geq 1$. Combining this with (4.23) implies

$$p_v(b)\|t_v\|_V \leq Cq_v \quad \text{where } q_v := \prod_{j \leq J_0} \eta^{v_j} \prod_{j > J_0} \frac{(1+v_j)^b}{\rho_j^{v_j}}, \quad (4.24)$$

with $\eta := \frac{1+\kappa}{2\kappa} < 1$ and $C := C_{\frac{r}{2}} C_0^{J_0}$. Using the inequality $1+x \leq e^x$ for $x > 0$, we deduce that for $j > J_0$ with $v_j \neq 0$

$$\frac{(1+v_j)^b}{\rho_j^{v_j}} \leq \left(\frac{|v_F|d_j}{v_j} \right)^{v_j} \quad \text{with } d_j := \frac{4B\|\psi_j\|_{L^\infty(D)}}{r}, \quad (4.25)$$

so we write $p_v(b)\|t_v\|_V \leq Cq_v \leq C\alpha(v_E)\beta(v_F)$, where

$$\alpha(v_E) := \prod_{j \leq J_0} \eta^{v_j}, \quad \beta(v_F) := \prod_{j > J_0} \left(\frac{|v_F|d_j}{v_j} \right)^{v_j}. \quad (4.26)$$

By (4.21) we find that $\sum_{j>J_0} d_j \leq \frac{1}{3}$. We then obtain the ℓ^p summability of $(q_v)_{v \in \mathcal{F}}$ exactly as in Sect. 3.2 in [13]. It follows that $(p_v(b)\|t_v\|_V)_{v \in \mathcal{F}}$ is ℓ^p summable. \square

The ℓ^p summability of $(p_v(b)\|t_v\|_V)_{v \in \mathcal{F}}$ implies by a standard result of Stechkin (see (3.13) in [13]) that there exist $C > 0$ and a sequence $(\Lambda_n)_{n \geq 1}$ of sets $\Lambda_n \subset \mathcal{F}$ of cardinality not exceeding n such that

$$\sum_{v \notin \Lambda_n} p_v(b)\|t_v\|_V \leq Cn^{-s}, \quad s := \frac{1}{p} - 1. \quad (4.27)$$

In this result, the sets Λ_n are built by picking n indices corresponding to the largest values of $p_v(b)\|t_v\|_V$ (note that these sets are not uniquely defined in the case where the same value is attained by several indices). However, these sets are not necessarily monotone. In order to obtain a sequence of monotone sets, we observe that the sequence $(q_v)_{v \in \mathcal{F}}$ defined in (4.24) is monotone decreasing, in the sense that $\mu \leq v \Rightarrow q_v \leq q_\mu$. Indeed, on the one hand, $q_{v+e_j} = \eta q_v < q_v$ if $j \leq J_0$. On the other hand, since $(2+n)^b \leq 2^b(1+n)^b \leq B(1+n)^b$ for any $n \geq 0$ and since $\frac{v_j+1}{|v_F|+1} \geq \frac{v_j}{|v_F|}$ for any $j > J_0$, we have

$$\frac{(2+v_j)^b}{\left(B + \frac{r(v_j+1)}{4(|v_F|+1)\|\psi_j\|_{L^\infty(D)}}\right)^{v_j+1}} \leq \frac{(1+v_j)^b}{\left(B + \frac{rv_j}{4|v_F|\|\psi_j\|_{L^\infty(D)}}\right)^{v_j}}$$

which implies also that $q_{v+e_j} \leq q_v$ when $j > J_0$. This shows that the sequence $(q_v)_{v \in \mathcal{F}}$ is *monotone decreasing* (i.e. monotonically decreasing with respect to the semiordeering \leq of multi-indices in \mathcal{F} induced by monotonicity in the sense of Definition 1.13).

Since the sequence $(q_v)_{v \in \mathcal{F}}$ majorises the sequence $(p_v\|t_v\|_V)_{v \in \mathcal{F}}$ up to a multiplicative constant, and since its ℓ^p summability is established in the above theorem, we also obtain (4.27) with Λ_n a set of indices corresponding to the n largest values of q_v . Since $(q_v)_{v \in \mathcal{F}}$ is monotone decreasing, such sets Λ_n can always be chosen to be monotone and nested. Combining these observations with Lemma 4.2, we thus have established the following convergence result.

Theorem 4.4 *Assume that if the diffusion coefficient satisfies the uniform ellipticity assumption (1.8) and one has $(\|\psi_j\|_{L^\infty(D)})_{j \geq 1} \in \ell^p(\mathbb{N})$ for some $0 < p < 1$. If the univariate sequence $(z_k)_{k \geq 0}$ is chosen so that $\lambda_k \leq (k+1)^\theta$ for some $\theta > 0$, then there exists a sequence $(\Lambda_n)_{n \geq 1}$ of monotone sets Λ_n such that $\#(\Lambda_n) = n$ and*

$$\|u - I_{\Lambda_n} u\|_{L^\infty(\mathcal{P}, V)} \leq Cn^{-s}, \quad s = \frac{1}{p} - 1. \quad (4.28)$$

Remark 4.5 The above convergence result exploits the fact that the Taylor series itself converges, due to the holomorphic properties of the solution map. For other type of parametric model problems that do not have such properties, the Taylor series may fail to converge even if the parameter dependence is smooth. On the other hand, other types of polynomial approximation may still converge with a high rate, and therefore

so may polynomial interpolation, provided that the Lebesgue constant does not grow too fast. For such problems, we thus expect that our high-dimensional interpolation algorithm behaves well, provided that the sets Λ_n are properly selected, either by a-priori estimates or by an adaptive strategy such as described in Sect. 2.5.

5 Numerical Experiments

5.1 Scalar Valued Functions

We first consider the interpolation of scalar valued functions $u : \mathcal{P} \rightarrow \mathbb{R}$ where $\mathcal{P} = P^d$ with $P = [-1, 1]$. Our objective is to test the adaptive algorithm AI proposed in Sect. 2.5 in various ways:

- Ability to select good monotone index sets Λ_n , in particular when the function has anisotropic dependence on the variable.
- The effect of the choice of the univariate sequence $(z_k)_{k \geq 0}$, in particular on the robustness of the interpolation with respect to noise in the measurements.
- Robustness of the performance with respect to the variable dimension d when the function depends only of a few unknown variables, or when the dependence with respect to the variables is sufficiently anisotropic.

We consider three possible choices for the univariate sequence $(z_k)_{k \geq 0}$:

- Uniform sequence (U): $z_0 = 1, z_1 = -1, z_2 = 0$ and for $k > 1$ we set $z_{2k+1} = \frac{1}{2} \sum_{j=0}^n \varepsilon_j 2^{-j}$ where $k = \sum_{j=0}^n \varepsilon_j 2^j$ is the binary expansion of k and $z_{2k+2} = -z_{2k+1}$. Such a choice produces a uniform subdivision of $[-1, 1]$ of step size 2^{-j} for the particular sections (z_0, \dots, z_{2^j}) , and avoids accumulation of points on a region of the interval for the intermediate sections $(z_0, \dots, z_k), 2^j < k < 2^{j+1}$.
- Leja sequence (L): $z_0 = 1$ and the sequence z_k is defined recursively on $[-1, 1]$ by (2.25).
- R-Leja sequence (R): this is the projection on $[-1, 1]$ of a Leja sequence for the complex unit disk initiated at 1. The R-Leja sequence has an explicit structure which is very similar to that of the sequence U in the sense that (z_0, \dots, z_{2^j}) are Clenshaw–Curtis points, that is, the projections on the real axis of a uniform subdivision of the upper half-circle with end-points at -1 and 1 , see [8–10].

Our first example is the function of $d = 16$ variables

$$u_1(\mathbf{y}) = u(y_1, \dots, y_{16}) = y_3 \sin(y_4 + y_{16}), \quad (5.1)$$

that in fact only depends on three variables. Figure 1 displays the interpolation error in terms of $n = \#(\Lambda_n)$ for the AI algorithm based on three possible choices U, L, R, of univariate sequence $(z_k)_{k \geq 0}$.

We observe that the error decays, after a (short) preasymptotic phase (required to identify the set $\{3, 4, 16\}$ of “active” dimensions) reaches machine precision when the AI algorithm is based on the choices L and R for the univariate sequence, but not with the choice U, although u_1 is analytic over \mathbb{R}^d . Inspection of the index sets Λ_n generated by the algorithm for the three choices reveals that for $n = 10^3$, the

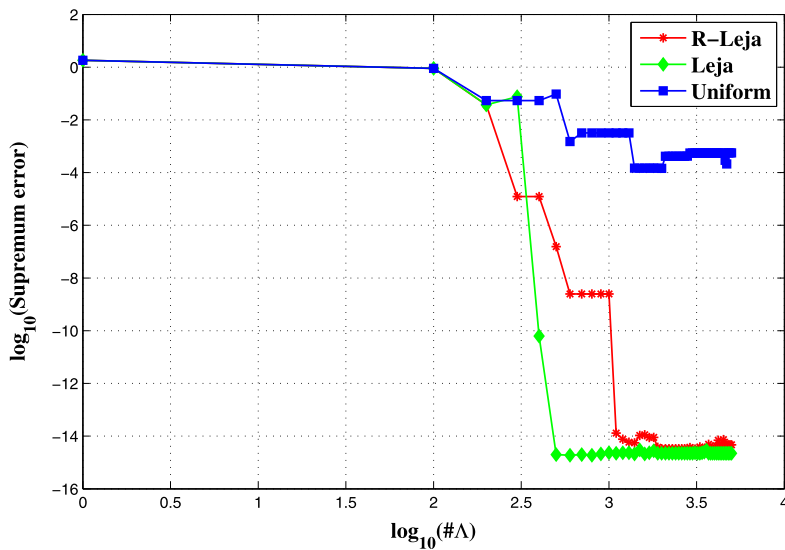


Fig. 1 Estimated L^∞ error of the AI algorithm applied to u_1 based on the sequences R , L and U

polynomial degree in the active variables y_4 and y_{16} reaches values above 30. This explains the bad behaviour of the error for the choice U . Indeed, it is well known that the univariate Lebesgue constant associated to k uniformly spaced points in $(-1, 1)$ grows faster than 2^k (see, e.g. [14]).

It appears that, in the range of polynomial degrees k activated by the AI algorithm based on the sequence U , the Lebesgue constants become so large that the (inevitable) roundoff errors in the floating point evaluation of the function $u_2(\mathbf{y})$ in (5.1) are amplified to $O(1)$ contributions to the interpolant. This does not occur with the choices L and R for which the Lebesgue constant has moderate growth. For these sequences, the algorithm identifies the three active variables (y_3, y_4, y_{16}) in the sense that all chosen indices ν have $\nu_j = 0$ for $j \neq 3, 4, 16$.

In our next example, we consider the parametric scalar function

$$u_2(\mathbf{y}) = \left(1 + \sum_{j=1}^d \gamma_j y_j\right)^{-1}, \quad \gamma_j := \frac{3}{5j^3}. \quad (5.2)$$

This function now depends on all variables y_1, \dots, y_d but in a strongly anisotropic way due to the decay of the weights γ_j . Since $\sum_{j=1}^\infty \gamma_j \approx 0.72 < 1$, the function u admits an holomorphic continuation (that is, an extension which is holomorphic in each variable) in the complex polydisc domain

$$\mathcal{P}_{\mathbb{C}} = (P_{\mathbb{C}})^d, \quad (5.3)$$

where $P_{\mathbb{C}}$ is the unit disk of \mathbb{C} , regardless of the dimension d . Moreover, an analysis completely analogous to that used in [13] to prove (1.10) for parametric PDE's, based on Cauchy's integral formula, reveals that, since $(\gamma_j)_{j \geq 1} \in \ell^p(\mathbb{N})$ for $p > \frac{1}{3}$, there

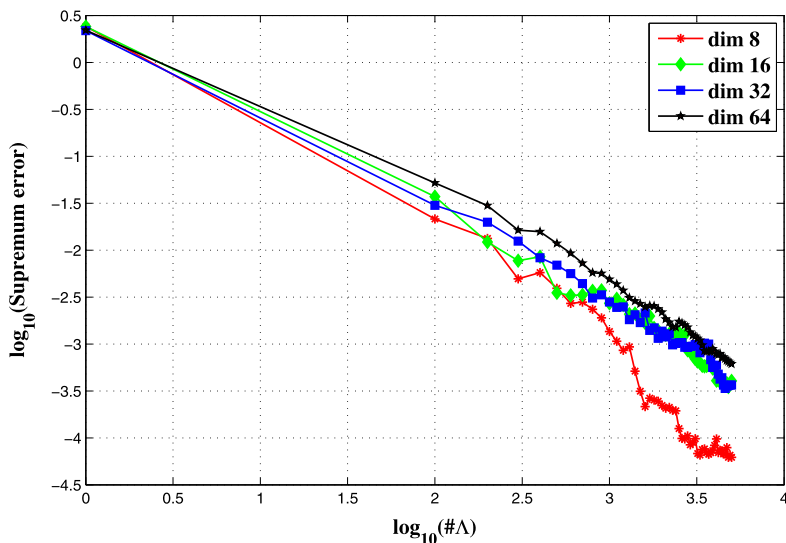


Fig. 2 Estimated L^∞ error for the AI algorithm applied to u_2 based on the sequences R for the dimensions $d = 8, 16, 32, 64$

exists a sequence $(\Lambda_n^*)_{n \geq 1}$ with $n = \#(\Lambda_n^*)$ such that

$$\inf_{v \in \mathbb{P}_{\Lambda_n^*}} \|u - v\|_{L^\infty(\mathcal{P})} \leq C_s n^{-s}, \quad n \geq 1, \quad (5.4)$$

for all $s < 2$ where C_s is independent of the number d of active coordinates. Figure 2 reveals that this robustness with respect to d is also observed when using the AI algorithm (here based on the sequence R), since its convergence behaviour is almost unchanged for $d = 8, 16, 32$ and 64 .

We next fix $d = 16$ and compare the error of the AI algorithm applied to u_2 based on the three sequences L , R and U . Figure 3 reveals that, in contrast to the function u_1 , the uniform sequence U gives as good results as the sequences L and R . This can be explained by inspecting more closely the index sets Λ_n , for which one finds that for $n = 10^4$ the highest polynomial degree attained on the most significant variable y_1 is 17 (due to the presence of many active variables) and the amplification of roundoff in function evaluations to $O(1)$ is not yet visible.

To verify that the preceding results and findings are not a consequence of the particular functional form of the exact solutions, we repeat the previous experiment with a small amplitude added random perturbation. Specifically, we now adaptively interpolate the values $u_2(\mathbf{z}_v) + \varepsilon_v$ where ε_v are independent realisations of a random variable with uniform law on $[-10^{-3}, 10^{-3}]$. Figure 4 reveals that the error diverges when using the U sequence, while it decays when using R or L (however, not reaching arbitrarily small values due to presence of the noise).

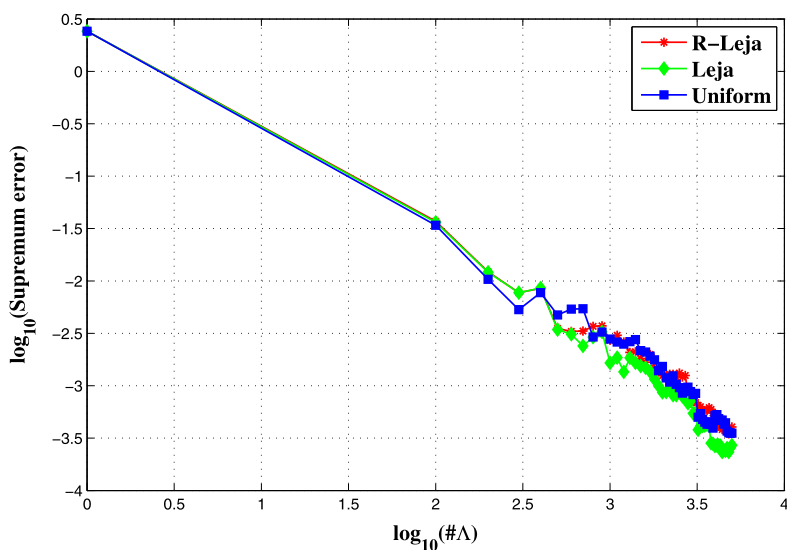


Fig. 3 Estimated L^∞ error of the AI algorithm applied to u_2 with $d = 16$ based on the sequences R , L and U

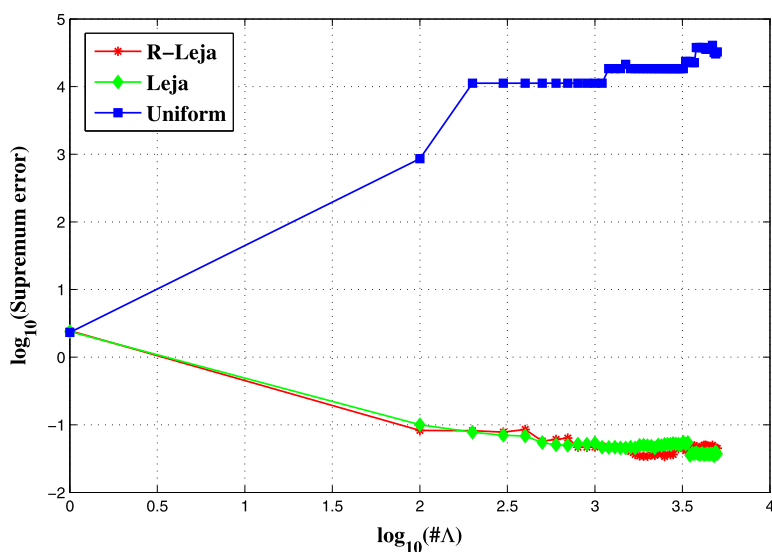


Fig. 4 Estimated L^∞ error of the AI algorithm applied to a noisy evaluation of u_2 with $d = 16$ based on the sequences R , L and U

Finally, we consider with $d = 16$ the function

$$u_3(\mathbf{y}) = \left(1 + \left(\sum_{j=1}^d \gamma_j y_j \right)^2 \right)^{-1}, \quad \gamma_j := \frac{5}{j^3}. \quad (5.5)$$

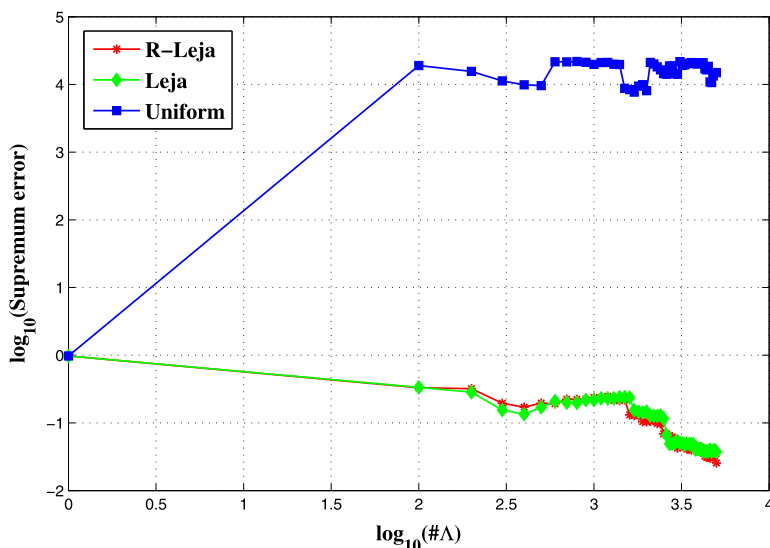


Fig. 5 Estimated L^∞ error of the AI algorithm applied to u_3 with $d = 16$ based on the sequences R , L and U

Similar to u_2 , this function has an anisotropic behaviour. It also has an holomorphic continuation in a complex neighbourhood of \mathcal{P} , however, it has smaller domain of analyticity than the polydisc $\mathcal{P}_{\mathbb{C}}$ due to the fact that $\sum_{j=1}^d \gamma_j > 1$. As a consequence, the AI algorithm based on the uniform sequence U does not converge *even in the noiseless case*, as illustrated by Fig. 5. This should be viewed as a manifestation of the well known Runge phenomenon.

In summary, the AI algorithm takes advantage of an anisotropic dependence on the variables, however, its success is critically tied to the choice of the univariate sequence $(z_k)_{k \geq 0}$ in either one of these situations: (i) the polynomial degree reaches high values in certain variables, (ii) the measurements (i.e. function evaluations) are “noisy”, (iii) the function does not have sufficient smoothness in certain variables. In all cases, both sequences R and L are good choices. The convergence rates are dimension independent if the function u allows for dimensionally independent best N -term approximation rates. The practical performance of the AI algorithms is found to be robust with respect to the choice of univariate sequence $(z_j)_{j \geq 1}$ as long as the sectional Lebesgue constants λ_k grow algebraically as k^θ with moderate exponent θ .

5.2 Parametric PDE's

We now turn to the interpolation of functions, $u : \mathcal{P} \rightarrow V$ defined as the solution map of a parametric PDE (1.1) where V is the solution space. In practice, the PDE is solved by a numerical technique such as the finite element method applied with a certain mesh, and therefore we rather interpolate the numerical solution map

$$u_h : \mathcal{P} \rightarrow V_h, \quad (5.6)$$

where $V_h \subset V$ is finite-dimensional and independent of the particular interpolation point in U .

In [11], several adaptive algorithms based on the Taylor partial sums (4.7) were proposed, analysed and implemented for the model elliptic PDE (1.7). The most practically efficient of these algorithms acts in a very similar way as the AI algorithm, in the sense that the set Λ_{n+1} is defined by adding to Λ_n the index v that maximises the V -norm of the Taylor coefficient t_v among the set of neighbours $\mathcal{N}(\Lambda_n)$. In the sequel, we refer to this adaptive algorithm as Largest Neighbour Taylor (LNT). Note that, in contrast to AI, algorithms based on the computation of the Taylor series such as LNT are by essence intrusive and strongly benefit from the particular structure of the problem (1.7): a linear equation with affine dependence of the operator on the parameters. This structure allows us to do a recursive computation of the Taylor coefficients t_v which are solutions to the boundary value problems

$$\int_D \bar{a} \nabla t_v \nabla v = - \sum_{j: v_j \neq 0} \int_D \psi_j \nabla t_{v-e_j} \nabla v, \quad v \in V. \quad (5.7)$$

These problems can be derived by partial differentiation at $\mathbf{y} = 0$ of the variational formulation

$$\int_D a(\mathbf{y}) \nabla u(\mathbf{y}) \nabla v = \int_D f v, \quad v \in V, \quad (5.8)$$

and we refer to [13] for a rigorous justification.

We compare the two algorithms AI and LNT when applied to (1.7) with $D = [0, 1]^2$ and diffusion coefficient $a(x, \mathbf{y})$ given by an expansion in the two-dimensional Haar wavelet basis, similar to Test 2 in [11], namely

$$a(x, \mathbf{y}) := \bar{a}(x) + \sum_{l=0}^L \beta_l \sum_{i=1}^3 \sum_{k \in \{0, \dots, 2^l-1\}^2} y_{l,k,i} h_{l,k}^i(x), \quad \bar{a} = 1. \quad (5.9)$$

In the above expansion,

$$h_{l,k}^i(x) := h^i(2^l x - k), \quad l \in \mathbb{N}, \quad k = (k_1, k_2) \in \{0, \dots, 2^l - 1\}^2, \quad i = 1, 2, 3, \quad (5.10)$$

where the generating wavelets h^i are defined by $h^1(x_1, x_2) := \varphi(x_1)h(x_2)$, $h^2(x_1, x_2) := h(x_1)\varphi(x_2)$ and $h^3(x_1, x_2) := h(x_1)h(x_2)$, with $\varphi := \chi_{[0,1]}$ and $h := \chi_{[0,1/2]} - \chi_{[1/2,1]}$. Using the relabelling

$$\psi_j := \beta_l h_{l,k}^i \quad \text{and} \quad y_j := y_{l,k,i}, \quad \text{when } j = 2^{2l} + 3(2^l k_1 + k_2) + i - 1, \quad (5.11)$$

we may rewrite the above expansion (5.9) in the form $a(x, y) := \bar{a}(x) + \sum_{j=1}^d y_j \psi_j(x)$ adopted in this paper, with $d := 2^{2(L+1)} - 1$. We consider the general form

$$\beta_l := c 2^{-\gamma l}, \quad c := 0.3(1 - 2^{-\gamma}), \quad (5.12)$$

which ensures that the uniform ellipticity assumption **UEA** holds with $r = 0.1$ and $R = 1.9$, regardless of the parametric dimension d . The value of the parameter $\gamma > 0$

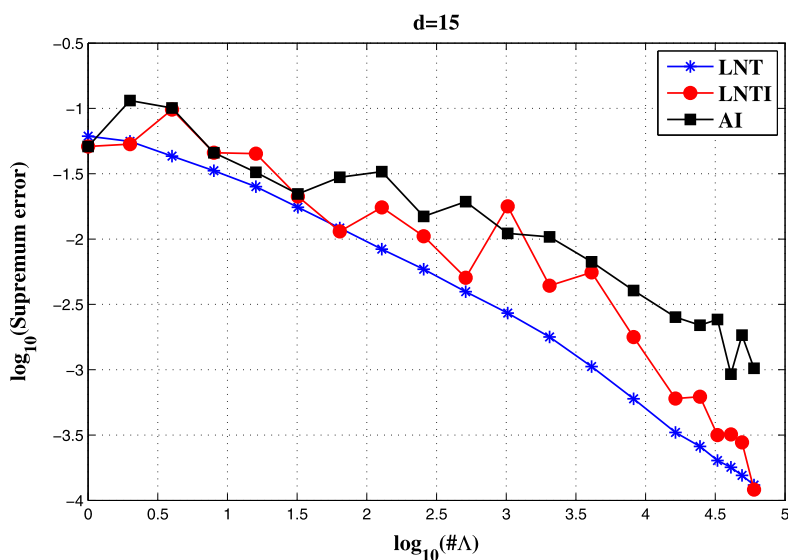


Fig. 6 Estimated $L^\infty(\mathcal{P}, V)$ error of LNT, LNTI and AI for the model problem (1.7) with coefficients (5.9) and $d = 15$

reflects the decay of the high scale oscillation and therefore the long range correlation in the diffusion field.

In our numerical test, we use the value $\gamma = 3$, which was among those tested in [11] and we consider the maximal scale levels $L = 1$ and 2 which give parametric dimension $d = 15$ and 63. In order to refine the comparison between AI and LNT, we introduce a third process that builds the interpolants $I_{\Lambda_n} u$ by using the sets Λ_n produced by the LNT algorithm. We refer to this algorithm as LNTI (Largest Neighbour Taylor Interpolation). For both AI and LNTI we use the univariate sequence R .

Several observations may be drawn from the error curves, displayed on Figs. 6 and 7. We first notice that the error curve of LNTI is close to that of LNT, with a more oscillatory behaviour. Since the sets $(\Lambda_n)_{n \geq 1}$ are the same for both algorithms, these oscillations may reflect the non-monotonic growth Lebesgue constant of the interpolation operator I_{Λ_n} that is used in LNTI. In addition, we notice that the error curve of AI is above that of LNTI, which means that AI is slightly misled in the adaptive selection of the sets Λ_n , which was better performed by LNT. In all cases, we find that these algorithms are robust with respect to the growth in the dimension, and are therefore capable of capturing the anisotropic structure of the problem reflected by the decay in the weights β_l .

The effectiveness of Taylor approximations in the above problem is due to the fact that the map $\mathbf{y} \mapsto u(\mathbf{y})$ has an holomorphic continuation in the complex polydisc $\mathcal{P}_{\mathbb{C}}$, similar to the scalar valued function u_2 . This approach becomes ineffective for PDE's where such continuation does not hold, even if the dependence on \mathbf{y} remains analytic over \mathcal{P} . Here are two instances of such problems:

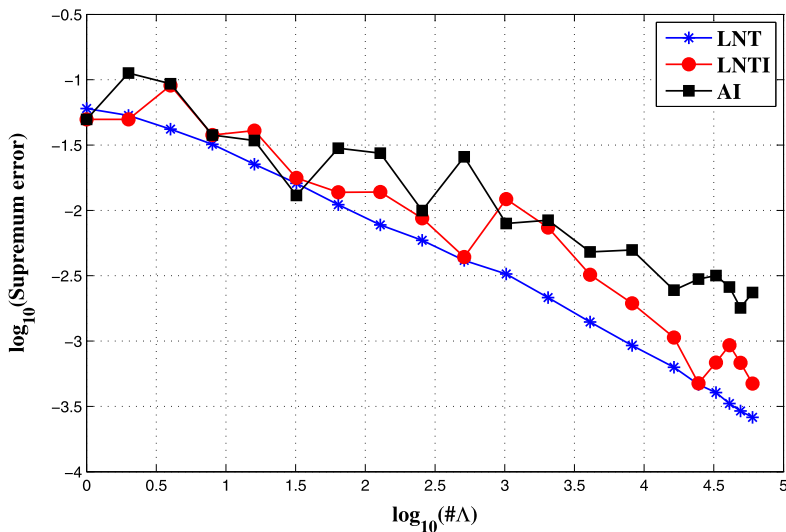


Fig. 7 Estimated $L^\infty(\mathcal{P}, V)$ error of LNT, LNTI and AI for the model problem (1.7) with coefficients (5.9) and $d = 63$

- Parametric nonlinear diffusion equations of the form

$$-\operatorname{div}(a \nabla u) + F(u) = f, \quad (5.13)$$

with similar assumptions on the diffusion coefficient a as in (1.7) and F an analytic nonlinear function. In the case where f is small enough, a perturbation approach leads to a similar holomorphic extension over the polydisc $\mathcal{P}_{\mathbb{C}}$ as in the linear case, see [20]. In the case where f is arbitrary, and F is such that the problem remains coercive in $H_0^1(D)$ (for example $F(u) = u^3$ in spatial dimension 2 or 3), analyticity still holds on \mathcal{P} , however, we are not ensured that the holomorphic extension can be carried over the whole of $\mathcal{P}_{\mathbb{C}}$.

- Linear PDE's set in parameter dependent domains. Consider for instance

$$-\Delta u = f \quad \text{in } D = D(\mathbf{y}) = \{(x_1, x_2) \mid 0 \leq x_1 \leq 1, 0 \leq x_2 \leq \phi(x_1, \mathbf{y})\}, \quad (5.14)$$

where the upper boundary has a general form $\phi(x_1, \mathbf{y}) = \bar{\phi}(x_1) + \sum_{j=1}^d y_j \phi_j(x_1) \geq \phi_{\min} > 0$. By a change of variable $x_2 \mapsto x_2 \phi(x_1)$ mapping the reference domain $\bar{D} = [0, 1]^2$ into $D(\mathbf{y})$, we are led to the elliptic PDE (1.7) over \bar{D} , however, with a matrix coefficient $a(x, \mathbf{y})$ depending in a rational manner of \mathbf{y} rather than affine. Here again, the dependence on \mathbf{y} is analytic over the real domain \mathcal{P} , however, with no guarantee that the holomorphic extension can be carried over the whole of $\mathcal{P}_{\mathbb{C}}$.

For such problems, we expect that Taylor-based algorithms such as LNT fail, while AI still converges. We illustrate this by studying a simpler problem, namely (1.7) with

diffusion coefficients of the form

$$a(x, \mathbf{y}) = \bar{a}(x) + \left(\sum_{j=1}^d y_j \psi_j(x) \right)^2. \quad (5.15)$$

Assuming that $\bar{a}(x) \geq r > 0$, the problem is always coercive and the map $\mathbf{y} \mapsto u(\mathbf{y})$ has analytic dependence on $\mathbf{y} \in \mathcal{P}$. However, when the function ψ_j are chosen too large, the holomorphic extension does not hold over the whole of $\mathcal{P}_{\mathbb{C}}$. This example should be viewed as similar to the scalar function u_3 .

Due to the polynomial form of $a(x, \mathbf{y})$ in the variable \mathbf{y} , it is still possible to compute the Taylor coefficients t_v in a recursive manner that generalises (5.7). Indeed, if $a(x, \mathbf{y})$ has the general form

$$a(x, \mathbf{y}) := \bar{a}(x) + \sum_{v \in \Lambda} \mathbf{y}^v \psi_v(x), \quad (5.16)$$

where Λ is a fixed set of index that does not include $0_{\mathcal{F}}$, then partial differentiation of (5.8) shows that

$$\int_D \bar{a} \nabla t_v \nabla v = - \sum_{\mu \in \Lambda : \mu \leq v} \int_D \psi_{\mu} \nabla t_{v-\mu} \nabla v, \quad v \in V. \quad (5.17)$$

This allows us to again perform LNT in the case where a has the quadratic form (5.15).

As an example, we take $d = 7$, $\bar{a}(x) = 1$ and

$$\psi_j(x) = \frac{\gamma}{j^3} \phi_j(x), \quad (5.18)$$

where $\gamma > 0$ and

$$\begin{aligned} \phi_1(x) &= 1, & \phi_2(x) &= \cos(2\pi x_1), & \phi_3(x) &= \cos(2\pi x_2), \\ \phi_4(x) &= \cos(2\pi(x_1 + x_2)), & \phi_5(x) &= \sin(2\pi x_1), \\ \phi_6(x) &= \sin(2\pi x_2), & \phi_7(x) &= \sin(2\pi(x_1 + x_2)). \end{aligned}$$

We compare LNT with AI, using both U and R univariate sequences for the second algorithm, as well as the LNTI algorithm that corresponds to use the interpolation algorithm (based on the R univariate sequence) but using the polynomial spaces generated by the LNT algorithm.

Figures 8 and 9 show the error curves, for both choices $\gamma = 0.5$ and $\gamma = 2$. When $\gamma = 0.5$, the LNT algorithm still converges due to the fact that the holomorphic extension holds over the whole of $\mathcal{P}_{\mathbb{C}}$, but it diverges for the larger value $\gamma = 2$. In contrast, the error still decreases when using AI based on the R univariate sequence. Similar to the scalar test case u_3 in (5.5), this algorithm diverges if we use the U sequence, due to the fast growth of the univariate Lebesgue constants λ_k for the U sequence. For both tests, we observe that the error of the AI algorithm decays and finally stagnate at some minimal error level. This level differs between the case $\gamma = 0.5$ and $\gamma = 2$,

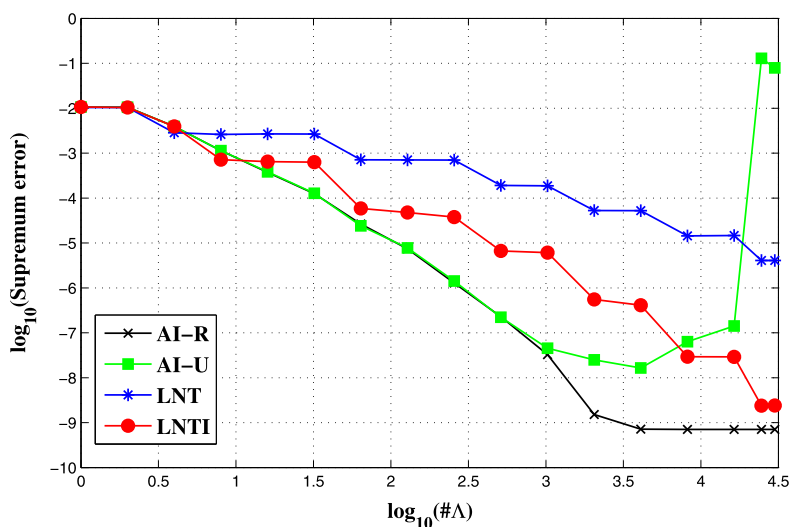


Fig. 8 Estimated $L^\infty(\mathcal{P}, V)$ error of LNT, of AI using U and R sequences, and of LNTI for the model problem (1.7) with coefficients (5.15) and $\gamma = 0.5$

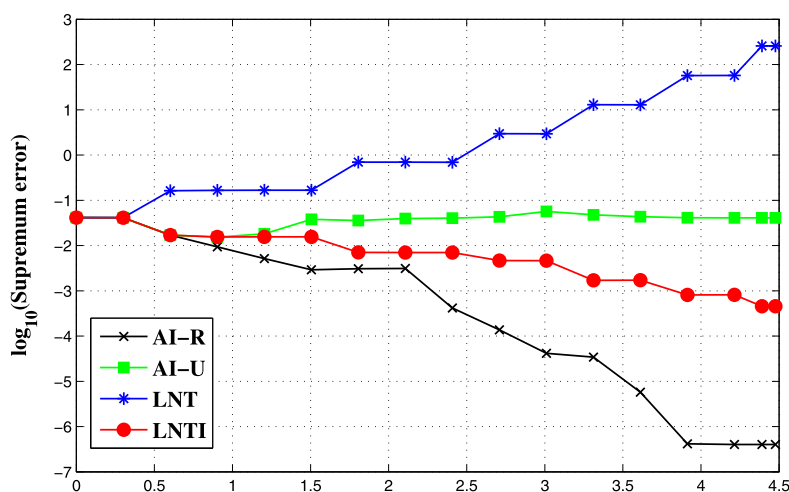


Fig. 9 Estimated $L^\infty(\mathcal{P}, V)$ error of LNT, of AI using U and R sequences, and of LNTI for the model problem (1.7) with coefficients (5.15) and $\gamma = 2$

reflecting the unavoidable dependence of the error with respect to the target function. Similar to the results displayed on Figs. 6 and 7 our interpretation is that the adaptive algorithm AI tends to be misled at some stage in the selection of the right polynomial spaces. Our detailed inspection of the index sets Λ_n generated by AI confirms this interpretation.

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